# So You Think You Can Divide? 

A History of Division

## Stephen Lucas

Department of Mathematics and Statistics James Madison University, Harrisonburg VA


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"There is a story of a German merchant of the fifteenth century, which I have not succeeded in authenticating, but it is so characteristic of the situation then existing that I cannot resist the temptation of telling it. It appears that the merchant had a son whom he desired to give an advanced commercial education. He appealed to a prominent professor of a university for advice as to where he should send his son. The reply was that if the mathematical curriculum of the young man was to be confined to adding and subtracting, he perhaps could obtain the instruction in a German university; but the art of multiplying and dividing, he continued, had been greatly developed in Italy, which in his opinion was the only country where such advanced instruction could be obtained."

## Outline

- Ancient Techniques
- Definitions
- Successive Subtraction
- Doubling
- Geometry
- Positional Notation
- Positional Definition
- Galley or Scratch
- Factor
- Napier's Rods and the "Modern" method
- Short Division and

Genaille's Rods

- Double Division
- Division Yielding Decimals
- Integer Division
- Modern Division
- Multiply by Reciprocal
- Iteration - Newton
- Iteration - Goldschmidt
- Iteration - EDSAC


## Definitions

If $a$ and $b$ are natural numbers and $a=q b+r$, where $q$ is a nonnegative integer and $r$ is an integer satisfying $0 \leq r<b$, then $q$ is the quotient and $r$ is the remainder after integer division. Also, $a$ is the dividend and $b$ is the divisor.

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$a / b$ takes a whole, divides it into $b$ parts, chooses $a$. Dividing each part into $d$ smaller parts means $a / b=(a d) /(b d)$. So $\frac{a / b}{c / d}=\frac{(a d) /(b d)}{(b c) /(b d)}=\frac{a d}{b c}$ in terms of the smaller pieces.

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Or, $100=0 \cdot 12+100=1 \cdot 12+88=2 \cdot 12+76=3 \cdot 12+64=$ $4 \cdot 12+52=5 \cdot 12+40=6 \cdot 12+28=7 \cdot 12+16=8 \cdot 12+4$.

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8184
16368
32736
$64 \quad 1472 \quad 1652-1472=180$
$1+2+4+64=71$,
so $1652 / 23=71$ r 19 .

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(b): $a>1$ so $a / 1=b / A D, A D=b / a$.

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Tens digit: $1652-70 \times 23=1652-1610=42$. Ones digit: $42-1 \times 23=42-23=19.1652 / 23=71 \mathrm{r} 19$ as before.
It would be nice to set out the computation more cleanly.

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\end{tabular}
e.g. \(9592 / 47\) : \(47 \times 1=47,47 \times 2=\)
\(94,47 \times 3=141,47 \times 4=188,47 \times\)
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Working backwards, $8095=(144 \times 8+4) \times 7+3=144 \times 8 \times 7$ $+4 \times 7+3=144 \times 56+31$,

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$24286=8095 \times 3+1,8095=1156 \times 7+3,1156=144 \times 8+4$.
Working backwards, $8095=(144 \times 8+4) \times 7+3=144 \times 8 \times 7$ $+4 \times 7+3=144 \times 56+31, \quad 24286=(144+31) \times 3+1$
$=144 \times 56 \times 3+31 \times 31+1=144 \times 168+94$.

## Factor Division

Fibonacci suggested splitting the divisor if possible:
$a \div(b c)=(a \div b) \div c$. Galley division is easier when the divisor is small, particularly single digits.
E.g. $24286 \div 168=24286 \div(3 \times 7 \times 8)$.
$24286=8095 \times 3+1,8095=1156 \times 7+3,1156=144 \times 8+4$.
Working backwards, $8095=(144 \times 8+4) \times 7+3=144 \times 8 \times 7$ $+4 \times 7+3=144 \times 56+31, \quad 24286=(144+31) \times 3+1$
$=144 \times 56 \times 3+31 \times 31+1=144 \times 168+94$.
But, finding factors means more division problems, factors may not be small, and multiple single digit divisions aren't easier than a single division.

## Napier's Rods

| 0 |
| :---: |
| $0 / 0$ |
| $0 / 0$ |
| $0 / 0$ |
| $0 / 0$ |
| $0 / 0$ |
| $0 / 0$ |
| 0 |
| 0 |
| 0 |
| 0 |
| 0 |
| 0 |




## Napier's Rods, Divisor Multiples, the Modern Method



## Napier's Rods, Divisor Multiples, the Modern Method



244 392/878:

$$
\left.\begin{array}{llllllll}
8 & 7 & 8
\end{array}\right) \begin{array}{llllll}
2 & 4 & 4 & 3 & 9 & 2
\end{array}
$$

## Napier's Rods, Divisor Multiples, the Modern Method



244 392/878:

$$
\left.\begin{array}{lll}
8 & 7 & 8
\end{array}\right)
$$

## Napier's Rods, Divisor Multiples, the Modern Method

| 1 |
| :---: |
| $0 / 2$ |
| $0$ |
| 0/4 |
| $0 / 5$ |
| $0 / 6$ |
| $0 / 7$ |
| $0$ |
| $0 / 9$ |



244 392/878:

$$
\left.\begin{array}{lllllll}
8 & 7 & 8
\end{array}\right) \begin{array}{llllll} 
& & & & \\
\hline 2 & 4 & 4 & 3 & 9 & 2 \\
1 & 7 & 5 & 6 & & \\
\hline & 6 & 8 & 7 & 9 &
\end{array}
$$

## Napier's Rods, Divisor Multiples, the Modern Method

| 1 |
| :---: |
| $0 / 2$ |
| $0$ |
| 0/4 |
| $0 / 5$ |
| $0 / 6$ |
| $0 / 7$ |
| $0$ |
| $0 / 9$ |



244 392/878:

$\left.\begin{array}{llllllll}8 & 7 & 8\end{array}\right)$|  |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |

## Napier's Rods, Divisor Multiples, the Modern Method

| 1 |
| :---: |
| $0 / 2$ |
| $0$ |
| 0/4 |
| $0 / 5$ |
| $0 / 6$ |
| $0 / 7$ |
| $0$ |
| $0 / 9$ |



244 392/878:

$$
\left.\begin{array}{lllllll}
8 & 7 & 8
\end{array}\right) \begin{array}{llllll} 
& & & & 2 & 7 \\
\hline
\end{array}
$$

## Napier's Rods, Divisor Multiples, the Modern Method

| 1 |
| :---: |
| $0 / 2$ |
| $0$ |
| 0/4 |
| $0 / 5$ |
| $0 / 6$ |
| $0 / 7$ |
| $0$ |
| $0 / 9$ |



244 392/878:

$$
\left.\begin{array}{lllllll}
8 & 7 & 8
\end{array}\right) \begin{array}{llllll} 
& & & & 2 & 7 \\
\hline
\end{array}
$$

## Napier's Rods, Divisor Multiples, the Modern Method

| 1 |
| :---: |
| $0 / 2$ |
| $0$ |
| 0/4 |
| $0 / 5$ |
| $0 / 6$ |
| $0 / 7$ |
| $0$ |
| $0 / 9$ |



244 392/878:


## Other Layouts

English speaking world, China,
Japan, India:
697
$7 \lcm{4883}$
428
$\frac{63}{5} 3$
$\frac{49}{4}$

## Other Layouts

English speaking world, China, Japan, India:

$$
\begin{array}{rr}
697 \\
\hline \begin{array}{r}
6883 \\
42 \\
68
\end{array} & \frac{42}{68} \\
\frac{63}{53} & \frac{49}{4} \\
\frac{49}{4} &
\end{array}
$$

## Much of Latin

America:

$$
4883 \div 7=697
$$

## Other Layouts

English speaking world, China, Japan, India:

$$
\begin{gathered}
697 \\
\frac{4883}{68}
\end{gathered}
$$

$$
\frac{63}{5} 3
$$

$$
\frac{49}{4}
$$

Much of Latin
America:

$$
\begin{aligned}
& 4883 \div 7=697 \\
& \begin{array}{r}
4297 \\
\hline 68 \\
\frac{43}{5} 3 \\
683 \\
\frac{49}{4}
\end{array} \quad 53 \\
& \\
& \hline 4
\end{aligned}
$$

Mexico:

## More Layouts

Spain, Italy, France, Portugal, Romania, Russia:

$$
\begin{aligned}
& 4883 \mid 7 \\
& -42 \\
& \hline 68 \\
& -\frac{63}{6} 3 \\
& \hline-\frac{49}{4}
\end{aligned}
$$

Brazil and Colombia, no | before quotient.

## More Layouts

Spain, Italy, France, Portugal, Romania, Russia:

$$
\begin{aligned}
& \begin{array}{r}
48 \\
-43 \\
-42 \\
\hline 68 \\
-\frac{63}{5} 3 \\
\hline
\end{array} \\
& -\frac{49}{4}
\end{aligned}
$$

## France:

$$
\begin{array}{rrr}
48 & 8 & 7 \\
-42 & 697 \\
\hline 6 & 8 \\
-\frac{6}{4} 3 \\
\hline & 3 & \\
-\frac{49}{4} &
\end{array}
$$

No decimals.

Brazil and Colombia, no | before quotient.

## More Layouts

Spain, Italy, France, Portugal, Romania, Russia:

$$
\begin{aligned}
& 4883 \mid 7 \\
& -42 \\
& \hline 68 \\
& -\frac{63}{69} \\
& \hline-\frac{49}{4}
\end{aligned}
$$

Brazil and Colombia, no | before quotient.

France:


No decimals.

Germany, Norway, Poland, Croatia, Slovenia, Hungary, Czech Republic, Slovakia, Bulgaria:

$$
\begin{aligned}
& 4883: 7=697 \\
& -\quad 42 \\
& \hline 68 \\
& -\frac{63}{53} \\
& -\frac{49}{4}
\end{aligned}
$$

## Short Division

With single digit divisors, each step will involve at most two digit numbers, and subtraction leaves a one digit number.

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With single digit divisors, each step will involve at most two digit numbers, and subtraction leaves a one digit number. For compactness, carry the single digit to the left as a subscript - short division, as opposed to traditional long division.


## Genaille's Rods for Short Division

| -WN-O | のutwn-o | $u+\omega N-0$ | AWN-0 | W N-O | N-O | -0 | \% |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\infty$ | $\checkmark$ | の | $u$ | - | $\omega$ | N | $\bigcirc$ |



## Example: $4883 \div$ ?

|  | 4 |  | 8 |  | 8 | 3 | R | D |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $7$ |  |  |  |  | $\begin{aligned} & 1 \\ & 6 \end{aligned}$ |  | $\begin{array}{\|c\|} \hline 0 \\ 1 \end{array}$ | 2 |  |
| $\begin{aligned} & 1 \\ & 4 \\ & 8 \end{aligned}$ |  |  |  |  | $\left\{\begin{array}{l} 1 \\ 4 \\ 7 \end{array}\right.$ |  | $-\begin{aligned} & 0 \\ & 1 \\ & 2 \end{aligned}$ | 3 |  |
| $\begin{array}{\|l} 3 \\ 6 \\ 8 \end{array}$ |  |  |  |  | $\gg \begin{aligned} & 0 \\ & 3 \\ & 5 \\ & 8 \end{aligned}$ | $x$ | $\begin{array}{\|} \hline 0 \\ 1 \\ 2 \\ 3 \end{array}$ | 4 |  |
| $\begin{array}{\|l\|} \hline 0 \\ 2 \\ 4 \\ 6 \\ 8 \\ \hline \end{array}$ |  | 1 3 5 7 9 |  |  |  |  | 0 1 2 3 4 | 5 |  |

## Example: $4883 \div$ ?

|  | 4 |  | 8 |  | 8 |  | 3 | R | D |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{array}{\|l\|} \hline 2 \\ 7 \end{array}$ |  |  |  |  |  |  |  | 0 | 2 |  |
| $\begin{array}{\|l\|} \hline 1 \\ 4 \\ 8 \\ \hline \end{array}$ |  |  |  | $\begin{aligned} & 2 \\ & 6 \\ & 9 \end{aligned}$ |  |  |  | 0 | 3 |  |
| $\begin{array}{\|l\|} \hline 1 \\ 3 \\ 6 \\ 8 \\ \hline \end{array}$ |  |  |  |  |  |  |  | $\begin{aligned} & 0 \\ & 1 \\ & 2 \\ & 3 \end{aligned}$ | 4 |  |
| 0 <br> 2 <br> 4 <br> 6 <br> 8 |  | $\begin{aligned} & 1 \\ & 3 \\ & 5 \\ & 7 \\ & 9 \end{aligned}$ |  | 1 3 5 7 9 |  |  |  | 0 <br> 1 <br> 2 <br> 3 <br> 4 | 5 |  |

So
$4883 \div 2=2441$ r 1 ,

## Example: $4883 \div$ ?

So

$$
\begin{aligned}
& 4883 \div 2=2441 \text { r } 1, \\
& 4883 \div 3=1627 \text { r } 2,
\end{aligned}
$$

## Example: $4883 \div$ ?

|  | 4 |  | 8 |  | 8 |  | 3 | R | D |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{array}{\|l\|} \hline 2 \\ 7 \end{array}$ |  |  |  |  |  |  |  | 0 | 2 |  |
| $\begin{array}{\|l\|} \hline 1 \\ 4 \\ 8 \\ \hline \end{array}$ |  |  |  | $\begin{aligned} & 2 \\ & 6 \\ & 9 \end{aligned}$ |  |  |  | 0 | 3 |  |
| $\begin{array}{\|l\|} \hline 1 \\ 3 \\ 6 \\ 8 \\ \hline \end{array}$ |  |  |  |  |  |  |  | $\begin{aligned} & 0 \\ & 1 \\ & 2 \\ & 3 \end{aligned}$ | 4 |  |
| 0 <br> 2 <br> 4 <br> 6 <br> 8 |  | $\begin{aligned} & 1 \\ & 3 \\ & 5 \\ & 7 \\ & 9 \end{aligned}$ |  | 1 3 5 7 9 |  |  |  | 0 <br> 1 <br> 2 <br> 3 <br> 4 | 5 |  |

So

$$
\begin{aligned}
& 4883 \div 2=2441 \text { r } 1, \\
& 4883 \div 3=1627 \text { r } 2, \\
& 4883 \div 4=1220 \text { r } 3
\end{aligned}
$$

So You Think You Can Divide?

## Example: $4883 \div$ ?

|  | 4 |  | 8 |  | 8 |  | 3 | R | D |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $7$ |  |  |  |  |  |  |  | $-8$ | 2 |  |
| $\begin{array}{\|l\|} \hline 1 \\ 4 \\ 8 \end{array}$ |  |  |  | $\begin{aligned} & 2 \\ & 6 \\ & 9 \end{aligned}$ |  |  |  | $\left\{\begin{array}{l} 0 \\ 1 \\ 2 \end{array}\right.$ | 3 |  |
| $\begin{array}{\|l\|} \hline 1 \\ 3 \\ 6 \\ 8 \\ \hline \end{array}$ |  | $\begin{array}{\|r\|} \hline 2 \\ 4 \\ 7 \\ 9 \\ \hline \end{array}$ |  |  |  |  |  | $\left\{\begin{array}{l} 0 \\ 1 \\ 2 \\ 3 \end{array}\right.$ | 4 |  |
| $\begin{array}{\|l\|} \hline 0 \\ 2 \\ 4 \\ 6 \\ 8 \end{array}$ |  | $\begin{array}{\|l\|} \hline 1 \\ 3 \\ 5 \\ 7 \\ 9 \end{array}$ |  | $\begin{array}{\|l\|} \hline 1 \\ 3 \\ 5 \\ 7 \\ 9 \end{array}$ |  |  |  | 0 1 2 3 4 | 5 |  |

So

$$
\begin{aligned}
& 4883 \div 2=2441 \text { r } 1, \\
& 4883 \div 3=1627 \text { r } 2, \\
& 4883 \div 4=1220 \text { r } 3, \\
& 4883 \div 5=976 \text { r } 3, \\
& \text { and so on. }
\end{aligned}
$$

## Double Division

# Chunking (UK, late 1990's) takes away "easy" (100, 10, 5, 2 etc.) multiples of the divisor. 

## Double Division

> Chunking (UK, late 1990's) takes away "easy" (100, 10, 5, 2 etc.) multiples of the divisor. Double division (Jeff Wilson, 2005) uses the first three doubles of the divisor as the chunks.

## Double Division

$$
7 \quad 4 \begin{array}{llll}
4 & 8 & 8 & 3
\end{array}
$$

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## Double Division

| $1_{\times}$ | 7 | $)$ | 4 | 8 | 8 | 3 |  |
| :--- | ---: | :--- | ---: | :--- | :--- | :--- | :--- |
| $2_{\times}$ | 14 | 2 | 8 | 0 | 0 | 400 |  |
| $4_{\times}$ | 28 | 2 | 0 | 8 | 3 |  |  |
| $8_{\times}$ | 56 | 1 | 4 | 0 | 0 |  |  |

Chunking (UK, late 1990's) takes away "easy" (100, 10, 5, 2 etc.) multiples of the divisor. Double division (Jeff Wilson, 2005) uses the first three doubles of the divisor as the chunks.

## Double Division

| $1_{\times}$ | 7 | $)$ | 4 | 8 | 8 | 3 |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $2 \times$ | 14 | 2 | 8 | 0 | 0 | 400 |  |
| $4 \times$ | 28 | 2 | 0 | 8 | 3 |  |  |
| $4_{\times}$ | 56 | 1 | 4 | 0 | 0 | 200 |  |
|  |  | 6 | 8 | 3 |  |  |  |
|  |  | 5 | 6 | 0 | 80 |  |  |
|  |  |  | 1 | 2 | 3 |  |  |

Chunking (UK, late 1990's) takes away "easy" (100, 10, 5, 2 etc.) multiples of the divisor. Double division (Jeff Wilson, 2005) uses the first three doubles of the divisor as the chunks.

## Double Division

| $1_{\times}$ | 7 | 4 | 8 | 8 | 3 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \times$ | 14 | 2 | 8 | 0 | 0 | 400 |
| $4 \times$ | 28 | 2 | 0 | 8 | 3 |  |
| $8 \times$ | 56 | 1 | 4 | 0 | 0 | 200 |
|  |  |  | 6 | 8 | 3 |  |
|  |  |  | 5 | 6 | 0 | 80 |
|  |  |  | 1 | 2 | 3 |  |
|  |  |  |  | 7 | 0 | 10 |
|  |  |  |  | 5 | 3 |  |

Chunking (UK, late 1990's) takes away "easy" (100, 10, 5, 2 etc.) multiples of the divisor. Double division (Jeff Wilson, 2005) uses the first three doubles of the divisor as the chunks.

## Double Division

| $1 \times$ | 7 | ) 4 | 4 | 8 | 3 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \times$ | 14 | 2 | 8 | 0 | 0 | 400 |
| $4 \times$ | 28 | 2 | 0 | 8 | 3 |  |
| $8 \times$ | 56 | 1 | 4 | 0 | 0 | 200 |
|  |  |  | 6 | 8 | 3 |  |
|  |  |  | 5 | 6 | 0 | 80 |
|  |  |  | 1 | 2 | 3 |  |
|  |  |  |  | 7 | 0 | 10 |
|  |  |  |  | 5 | 3 |  |
|  |  |  |  | 2 | 8 | 4 |
|  |  |  |  | 2 | 5 |  |

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## Another Example

## $214) \longdiv { 7 } 3 4 8 8$

## Another Example

$1_{\times} \quad 214$ ) $7 \begin{array}{lllll} & 3 & 4 & 8 & 5\end{array}$
$2 \times 428$
$4 \times 856$
$8 \times 1712$

## Another Example

|  | $1 \times$ | 214 | $)$ | 7 | 3 | 4 | 8 | 5 |
| :--- | ---: | :--- | ---: | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |  |  |  |  |
| $2 \times$ | 428 | 4 | 2 | 8 | 0 | 0 |  | 200 |
| $4 \times$ | 856 | 3 | 0 | 6 | 8 | 5 |  |  |
| $8 \times$ | 1712 |  |  |  |  |  |  |  |

## Another Example

| $1_{\times}$ | 214 | $)$ | 7 | 3 | 4 | 8 | 5 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |  |  |  |  |
| $2 \times$ | 428 | 4 | 2 | 8 | 0 | 0 |  |
| $4 \times$ | 856 | 200 |  |  |  |  |  |
| $4_{\times}$ | 3 | 0 | 6 | 8 | 5 |  |  |
| $8_{\times}$ | 1712 | 2 | 1 | 4 | 0 | 0 | 100 |
|  |  |  | 9 | 2 | 8 | 5 |  |

## Another Example



## Another Example

| $1^{\times}$ | 214 | 7 | 3 | 4 | 8 |  | 200 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $2 \times$ | 428 | 4 | 2 | 8 | 0 |  |  |
| $4 \times$ | 856 | 3 | 0 | 6 | 8 |  |  |
| $8 \times$ | 1712 | 2 | 1 | 4 | 0 |  | 100 |
|  |  |  | 9 | 2 | 8 |  |  |
|  |  |  | 8 | 5 | 6 |  | 40 |
|  |  |  |  | 7 | 2 |  |  |
|  |  |  |  | 4 | 2 |  | 2 |
|  |  |  |  | 2 | 9 |  |  |

## Another Example



## Another Example



## Using Integer Division

Given that $(p<q) \frac{p}{q}=\frac{a-1}{10}+\frac{a-2}{10^{2}}+\frac{a-3}{10^{3}}+\cdots$, successively multiply by ten, stop if repeated or zero remainder.

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Given that $(p<q) \frac{p}{q}=\frac{a_{-1}}{10}+\frac{a_{-2}}{10^{2}}+\frac{a_{-3}}{10^{3}}+\cdots$, successively multiply by ten, stop if repeated or zero remainder.

For example, $\frac{11}{16}=\frac{a_{-1}}{10}+\frac{a_{-2}}{10^{2}}+\frac{a_{-3}}{10^{3}}+\cdots$.
Times 10: $10 \times \frac{11}{16}=\frac{110}{16}=6 \frac{7}{8}$

## Using Integer Division

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so $a_{-1}=6$ and $\frac{7}{8}=\frac{a_{-2}}{10}+\frac{a_{-3}}{10^{2}}+\frac{a_{-4}}{10^{3}}+\cdots$.

## Using Integer Division

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Times 10: $10 \times \frac{7}{8}=\frac{70}{8}=8 \frac{3}{4}$

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Times 10: $10 \times \frac{11}{16}=\frac{110}{16}=6 \frac{7}{8}=a_{-1}+\frac{a_{-2}}{10}+\frac{a_{-3}}{10^{2}}+\frac{a_{-4}}{10^{3}}+\cdots$, so $a_{-1}=6$ and $\frac{7}{8}=\frac{a_{-2}}{10}+\frac{a_{-3}}{10^{2}}+\frac{a_{-4}}{10^{3}}+\cdots$.

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Times 10: $10 \times \frac{7}{8}=\frac{70}{8}=8 \frac{3}{4}=a_{-2}+\frac{a_{-3}}{10}+\frac{a_{-4}}{10^{2}}+\frac{a_{-5}}{10^{3}}+\cdots$,
so $a_{-2}=8$ and $\frac{3}{4}=\frac{a_{-3}}{10}+\frac{a_{-4}}{10^{2}}+\frac{a_{-5}}{10^{3}}+\cdots$.

## Integer Division Example, continued

Times 10: $10 \times \frac{3}{4}=\frac{30}{4}=7 \frac{1}{2}$

## Integer Division Example, continued

Times 10: $10 \times \frac{3}{4}=\frac{30}{4}=7 \frac{1}{2}=a_{-3}+\frac{a_{-4}}{10}+\frac{a_{-5}}{10^{2}}+\frac{a_{-6}}{10^{3}}+\cdots$,

## Integer Division Example, continued

$$
\begin{aligned}
& \text { Times } 10: 10 \times \frac{3}{4}=\frac{30}{4}=7 \frac{1}{2}=a_{-3}+\frac{a_{-4}}{10}+\frac{a_{-5}}{10^{2}}+\frac{a_{-6}}{10^{3}}+\cdots, \\
& \text { so } a_{-3}=7 \text { and } \frac{1}{2}=\frac{a_{-4}}{10}+\frac{a_{-5}}{10^{2}}+\frac{a_{-6}}{10^{3}}+\cdots .
\end{aligned}
$$

## Integer Division Example, continued

Times 10: $10 \times \frac{3}{4}=\frac{30}{4}=7 \frac{1}{2}=a_{-3}+\frac{a_{-4}}{10}+\frac{a_{-5}}{10^{2}}+\frac{a_{-6}}{10^{3}}+\cdots$,
so $a_{-3}=7$ and $\frac{1}{2}=\frac{a_{-4}}{10}+\frac{a_{-5}}{10^{2}}+\frac{a_{-6}}{10^{3}}+\cdots$.
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Of course, process could also be periodic.

## Decimal Long Division

Each of the steps of the previous example $(11 / 16)$ is equivalent to one step in dividing 110000 by 16 , just shifted to deal with the position of the decimal. So we can use standard long division with a decimal point.

## Decimal Long Division

$$
\begin{aligned}
& 990) \frac{22.3161}{22093.0000 \ldots} \\
& \begin{array}{r}
1980 \\
2293 .
\end{array} \\
& \begin{array}{r}
1980 . \\
\hline 313.0
\end{array} \\
& \frac{297.0}{16.00} \\
& \frac{9.90}{6.100} \\
& \frac{5.940}{1600} \\
& \begin{array}{r}
990 \\
\hline 61
\end{array}
\end{aligned}
$$

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For example, 22093/990 = 22.31616161616....

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This technique was used by the Ancient Babylonians in base sixty which is a highly composite number, so most reciprocals have a finite radix sixty representation.

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For example $1 / 7$ with $x_{0}=0.2$,
$x_{1}=x_{0}\left(2-7 \times x_{0}\right)=0.2(2-7 \times 0.2)=0.12$,
$x_{2}=x_{1}\left(2-7 \times x_{1}\right)=0.12(2-7 \times 0.12)=0.1392$,
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Accuracy doubles at every step, given a close enough initial guess.

## Graphical Proof



If $x \approx 1 / a$ and we ignore the top left rectangle, then $(1 / x) \cdot(1 / a) \approx a x+2(1-a x), 1 / a \approx x(a x+2-2 a x)$, or $1 / a \approx x(2-a x)$.

## Goldschmidt's Iteration

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To calculate $p / q$, start with $y_{0}$ close to $1 / q$. Then $p_{1} / q_{1}=\left(p y_{0}\right) /\left(q y_{0}\right)$ where $q_{1}$ is close to one.

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Algorithm: given $p_{i}$ and $q_{i}$ where $q_{i} \approx 1$, let $y_{i}=2-q_{i}$ then $p_{i+1}=y_{i} p_{i}$ and $q_{i+1}=y_{i} q_{i}$.

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Identical to Newton with $p_{i}=x_{i-1}$ and $q_{i}=a x_{i-1}$. So why use it?
The two multiplications can be done in parallel, essentially doubling the speed!

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where $p_{2}=p_{1}\left(1+c_{1}\right)$ and $c_{2}=c_{1}^{2}$.
In general, $p_{i+1}=\left(1+c_{i}\right) p_{i}$ and $c_{i+1}=c_{i}^{2}$ with $c_{1}=1-q y_{0}$.

## Conclusion

So, just how would you like to divide now?

